

Link-State Routing for Compute Nodes within a Cluster

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Goal: Implement an open source dynamic routing protocol to increase reliability and maintainability while reducing resource overhead for high performance computing clusters.

Current Configuration Overview

IO Nodes

- Gateways for compute nodes used to access high speed parallel storage file systems that contain crucial data required for scientific research.
- Failure of an IO node hinders data transmission leading to possible non-recoverable data loss and job failure.

Dead Gateway Detection

- Custom script developed and maintained by LANL to determine the current health state of IO nodes.
- ICMP (ping) is utilized to identify broken links
- Reconfigures routing information within each of the



Open Shortest Path First (OSPF)

- Interior gateway routing protocol utilizing the link-state routing algorithm.
- Utilizes a tree topology requiring all areas to be attached to the backbone area 0.
- Routers assign a cost to each network segment/link, lower integers identify a link with higher preference.
- A Route is calculated all at once by each router using Dijkstra's Shortest Path First (SPF) algorithm and added to the shortest-path tree.
- OSPF Hello packet (equivalent to a half ICMP) is utilized to keep neighbor adjacencies
- Network traffic is minimized by only updating the

routing information thru a process referred to as "reliable flooding".

Stretch Cluster

Designs Explored

Stretch Cluster: Stub vs Totally Stub Areas

	Summary Link S	tates (Area 0.0.0.0)
Link ID 172.16.0.0 192.168.0.0	ADV Router 10.15.6.1 10.15.6.1	Age Seq# CkSum Route 194 0x80000031 0xe65f 172.16.0.0/16 194 0x80000032 0xb8df 192.168.0.0/16
	Summary Link	States (Area 0.0.0.6 [Stub])
Link ID 0.0.0.0 10.15.1.0 10.15.6.0 172.16.0.0	ADV Router 10.15.6.1 10.15.6.1 10.15.6.1 10.15.6.1	Age Seq# CkSum Route 194 0x80000002 0x8eae 0.0.0.0/0 194 0x80000001 0x0d04 10.15.1.0/24 104 0x80000002 0x6fa5 10.15.6.0/24 194 0x80000002 0x6314 172.16.0.0/16
	Summary Link S	tates (Area 0.0.0.7 [Stub])
Link ID 0.0.0.0	ADV Router 10.15.6.1 AS External Li	Age Seq# CkSum Route 204 0x80000001 0x90ad 0.0.0.0/0 .nk States
Link ID 0.0.0.0	ADV Router 10.15.1.254	Age Seq# CkSum Route 1356 0x8000002f 0xda8f E2 0.0.0.0/0 [0x0]

Figure 2: Type 3 LSA comparison between areas

- OSPF Totally Stub areas prevent communication between clusters.
- Totally stub areas reduce inter-area network traffic.

Efforts to minimize production configuration changes displayed in *figure 1*, involved planning for an extra layer of separation for the compute nodes.

Compute Node Area Separation

- Virtual Links between areas not connected to the backbone and area 0 *Limitations:* Virtual links are not supported by Arista hardware & Quagga
- IS-IS (intermediate system to intermediate system) is a link-state routing protocol similar to OSPF supports areas in the form of levels.

Limitations: Redistribution and IP over IB not supported

OSPF Routing Domain Re-design

- Inner Router (IR): Designated Router (DR) Compute Nodes
- Area Border Router (ABR): IO Nodes
- Area X totally stub prevents communication between clusters
- Non-Broadcast network for compute nodes establishes only adjacency with IO nodes
- Broadcast configuration for IO nodes establishes adjacencies with all directly connected routers.
- Default-Routes defined by OSPF (allows external domain communication)



Figure 3: NEW OSPF Routing Domain

Results & Future Tests



Figure 4: TCP packet capture on Stretch IO node

Results & Observations:

- Compute nodes only form adjacencies with IO Nodes
- OSPF Total packet count is highest when service is started or restarted: Database description packets account for this spike
- Packet count is minimal with packet no larger than ~2,000 bits
- "Reliable Flooding" occurs every ~30 min

Future Tests:

- Kit Cluster: 4-8 IO Nodes ~30 Compute Nodes
- Moon Cluster: ~50 IO Nodes ~1500 Compute Nodes
 DGD overhead comparison to OSPF overhead
 OSPF protocol scalability Maximum Packet size ~ 65,000 Bytes



AcknowledgementsReferencesLANL HPC Networking Team, Lowell Wofford, Paul Peltz, Susan Coulter, Dave Morton, Hunter Easterday[1] Moy, J. T. (1998). OSPF: Anatomy of an Internet routing protocol. Reading, MA: Addison-Wesley.
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